- 1. 1. Batch, cookies are made according to a recipe.
 - 2. Discrete, discrete pieces are put together.
 - 3. Continuous, paper is continuously coming out of the paper machine.
 - 4. Continuous, electricity is continuously produced.
 - 5. Batch, medicine is produced according to a recipe.
 - 6. Batch, beverages are produced according to a recipe.
- **2.** See figure 1.
- 3.
 - **a.** Each inequality corresponds to a line in figure 2. The area in the bottom left of the figure is the feasible area.
 - b. The optimal value is always in one of the corners of the feasible area, i.e.

$$v_0 = \begin{pmatrix} 0\\0 \end{pmatrix}, v_1 = \begin{pmatrix} 1\\0 \end{pmatrix}, v_2 = \begin{pmatrix} 0\\2 \end{pmatrix}, v_3 = \begin{pmatrix} 2&2\\3&1 \end{pmatrix}^{-1} \begin{pmatrix} 4\\3 \end{pmatrix} = \begin{pmatrix} 0.5\\1.5 \end{pmatrix}$$

The values in the corners are $v_0: 0, v_1: 3, v_2: 2, v_3: 3$, meaning that e.g. v_1 or v_3 is optimal.

Alternatively one can note that the boundary $3x_1 + x_2 \leq 3$ is included in the feasible area, and thus the optimal value is 3 for any (x_1, x_2) on this line in the first quadrant.

- 4. a) single-path, single product only one predefined sequential path from raw materials to the only product A
 - **b)** multi-path, multi-grade three parallel branches without interaction from raw materials to the two variants of product A
 - c) network, single product interacting parallel branches only producing product A
 - **d)** multi-path, multi-product parallel branches without interaction, more than one product
 - e) network, single product interacting parallel branches only producing product A
 - **f)** multi-path, multi-product parallel branches without interaction, more than one product

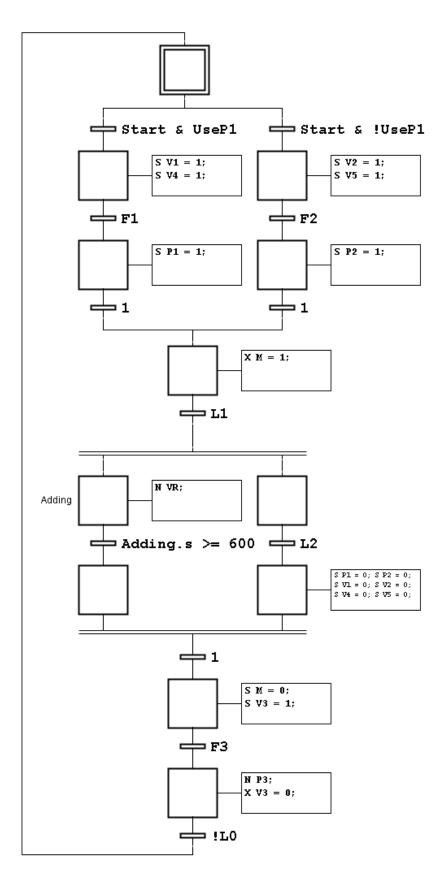


Figure 1 Grafcet for solution preparation process.

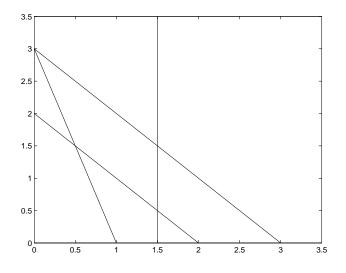


Figure 2 Feasible area.

5. No solution

6.

a. The area dependence matrix A_d gives a representation of the site structure and can be calculated from the flowchart of the product flow.

$$A_d = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix}$$

The area-utility matrix A_u gives a representation of which utilities are required in each area. Considering that the utilities are ordered: electricity, instrument air, and vacuum system; the area-utility matrix A_u becomes:

$$A_u = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix}$$

The utility dependence matrix U_d defines the interdependence between the different utilities. Considering the same ordering for the utilities as for A_u , then:

$$U_d = \begin{pmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{pmatrix}$$

The utility operation matrix U describes which utilities have operated correctly at each sample point.

$$U = \begin{pmatrix} 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 & 1 & 1 \end{pmatrix}$$

b. The total revenue loss due to each utility can be calculated using:

$$\begin{aligned} J_{u}^{tot} &= \operatorname{diag}(1 - U_{av}^{ud}) \cdot \operatorname{sign}(A_{d}A_{u})^{T}(q^{m} \cdot *p)n_{s}t_{s} \\ &= \operatorname{diag}\begin{pmatrix} 1/6\\1/6\\0 \end{pmatrix} \operatorname{sign}\begin{pmatrix} 1 & 1 & 3 & 2\\1 & 0 & 2 & 0\\1 & 1 & 2 & 1 \end{pmatrix} (q^{m} \cdot *p)n_{s}t_{s} \\ &= \begin{pmatrix} 1/6 & 0 & 0\\0 & 1/6 & 0\\0 & 0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 1 & 1 & 1\\1 & 0 & 1 & 0\\1 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} 2 \cdot 3\\0 \cdot 1\\4 \cdot 6\\5 \cdot 2 \end{pmatrix} \frac{\mathrm{m}^{3}}{\mathrm{h}} \cdot \frac{\mathrm{kr}}{\mathrm{m}^{3}} \cdot 6 \cdot \frac{1}{5}\mathrm{h} \\ &= \begin{pmatrix} 8 & 6 & 0 \end{pmatrix}^{T} \mathrm{kr} \end{aligned}$$

The utility that causes the greatest losses at the site is electricity.

7. Denote the number of rings and necklaces to produce x_1 and x_2 respectively and the utilization of bender 1 and 2 by x_3 and x_4 respectively. The optimization problem is then

maximize
$$15x_1 + 103x_2 - 100 - \frac{1}{2}x_4$$

subject to $x_3 \le 200$
 $x_4 \le 400$
 $4x_1 + 21x_2 \le x_3 + x_4$
 $x \ge 0$

8.

a. In this solution we consider the more general case where we have N companies.

We have the profit function for company i: $u_i = q_i(a-c-Q)$ where $Q = \sum_i q_i$. We find the Nash Equilibrium by differentiating: $\frac{\partial u_i}{\partial q_i} = a - c - Q - q_i = 0$ Summing these equations we get

$$0 = \sum_{i} a - c - Q - q_{i}$$

$$0 = N(a - c - Q) - Q$$

$$(N+1)Q = N(a - c)$$

$$Q = \frac{N}{N+1}(a - c)$$

Using $q_1^* = q_2^* = ... = q_N^*$ we get the Nash Equilibrium $q_i^* = \frac{Q}{N} = \frac{1}{N+1}(a-c)$. Thus for N = 3 we have $q_1^* = q_2^* = q_3^* = \frac{1}{4}(a-c)$.

b. The profit for company i is

 $u_i^* = \frac{1}{N+1}(a-c)\left(a-c-\frac{N}{N+1}(a-c)\right) = \frac{1}{N+1}(a-c)\left((a-c)\frac{1}{N+1}\right) = \frac{(a-c)^2}{(N+1)^2}.$ The total profit for all companies is $\frac{(a-c)^2}{(N+1)^2}N$, i.e. $(a-c)^2\frac{2}{9}$ and $(a-c)^2\frac{3}{16}$ for N = 2 and N = 3 respectively. Thus the total profit is smaller in the triopoly market.

9.

a. The following table describes the strategies of Player 1 and Player 2:

Player 2

$$1 \quad 2$$

Player 1 $1 \quad -2 \quad +3$
 $2 \quad +3 \quad -4$

- **b.** Analysing the game from the point of view of Player 1, suppose he calls "one" 3/5ths of the time and "two" 2/5ths of the time at random. In this case:
 - If Player 2 calls "one", Player 1 loses 2 SEK 3/5ths of the time and wins 3 SEK 2/5ths of the time, that is, on the average, he wins: -2(3/5) + 3(2/5) = 0 SEK (he breaks even in the long run).
 - If Player 2 calls 'two', Player 1 wins 3 SEK 3/5ths of the time and loses 4 SEK 2/5ths of the time, that is, on the average he wins: 3(3/5) 4(2/5) = 1/5 SEK.

Therefore if Player 1 mixes his choices in the given way, the game is even every time Player 2 calls "one", but Player 1 wins 1/5 SEK on the average every time Player 2 calls "two". In this way Player 1 is assured of at least breaking even on the average no matter what Player 2 does.

c. Considering "p" as the proportion of times that Player 1 calls "one", the idea is to choose "p" such that Player 1 wins the same amount on the average whether Player 2 calls "one" or "two". Then:

Player 2

$$1 2$$

Player 1 $p 1 -2 +3$
 $(1-p) 2 +3 -4$

The payoff of Player 1 if Player 2 calls "one" is: -2p + 3(1-p)The payoff of Player 1 if Player 2 calls "two" is: 3p - 4(1-p)Therefore Player 1 should choose "p" so that:

$$-2p + 3(1-p) = 3p - 4(1-p)$$

 $p = 7/12$

Thus, Player 1 should call "one" with probability 7/12, and "two" with probability 5/12. Using this strategy Player 1 wins 1/12 SEK every time he plays the game, no matter what Player 2 does.