Updated solutions to Problems 2.4, 4.1, and 6.1



Figure 2.1 Computer-controlled system.

- **2.4** Consider the computer-controlled system in Fig. 2.1. The P-controller should execute with the sampling interval h = 0.5 s, and the controller gain is given by K = 2.
 - a. Consider the following implementation of the controller:

```
LOOP
  y = readInput();
  u = -K*y;
  writeOutput(u);
  waitForNextPeriod();
END;
```

Assume that the execution time of the controller is negligible. Show that the closed-loop system is stable.

b. Now consider the following implementation of the controller:

```
LOOP
writeOutput(u);
y = readInput();
u = -K*y;
waitForNextPeriod();
END;
```

Will the closed-loop system still be stable?

c. Consider again the implementation from subproblem a, i.e.:

```
LOOP
  y = readInput();
  u = -K*y;
  writeOutput(u);
  waitForNextPeriod();
END;
```

Assume now that the execution time of the controller is not negligible, i.e. that there is a delay, L, between readInput() and writeOutput(u). Assuming that L is constant and L < h, what is the largest value of L for the system to be stable?

Hint: The stability conditions for a second-order discrete-time system with the characteristic polynomial $A(z) = z^2 + a_1 z + a_2$ are given by

$$a_2 < 1$$

 $a_2 > -1 + a_1$
 $a_2 > -1 - a_1$

Solution

a. Sampling the process using the table "Zero-order hold sampling of a continuoustime system with transfer function G(s)" gives

$$H(z) = \frac{0.6321}{z - 0.3679}$$

The closed-loop system is given by

$$H_{cl}(z) = rac{KH(z)}{1+KH(z)} = rac{1.264}{z+0.8964}$$

The pole is located in -0.8964, inside the unit circle, so the closed-loop system is stable.

b. The sampled process, including a one sample delay, is now given by

$$H(z) = \frac{0.6321}{z(z - 0.3679)}$$

The closed-loop system is given by

$$H_{cl}(z) = \frac{1.264}{z^2 - 0.3679z + 1.264}$$

The poles are located in $0.1836 \pm 1.1092i$, i.e., outside the unit circle, so the closed-loop system in unstable.

c. We start by writing the continuous-time system on state-space form, i.e.,

$$\frac{dx(t)/dt = -2x(t) + 2u(t)}{y(t) = x(t)}$$

The computational delay is equivalent to a constant input delay, i.e., the continuous-time system will be

$$\frac{dx(t)}{dt} = -2x(t) + 2u(t - L)$$
$$y(t) = x(t)$$

The ZOH-sampled equivalent of this, assuming that $L \leq h$ is

$$x(kh+h) = \Phi x(kh) + \Gamma_0 u(kh) + \Gamma_1 u(kh-h)$$

$$y(kh) = x(kh)$$

where

$$\begin{split} \Phi &= e^{-2h} = e^{-1} \\ \Gamma_0 &= 2 \int_0^{h-L} e^{-2s} ds = 1 - e^{2L-1} \\ \Gamma_1 &= 2e^{-2(h-L)} \int_0^L e^{-2s} ds = e^{2L-1} - e^{-1} \end{split}$$

Applying the control law u(k) = -2y(k) = -2x(k) gives the closed loop system

$$x(k+1) = e^{-1}x(k) - 2(1 - e^{2L-1})x(k) - 2(e^{2L-1} - e^{-1})x(k-1)$$

The characteristic equation is hence

$$z^{2} + (2(1 - e^{2L-1}) - e^{-1})z + 2(e^{2L-1} - e^{-1})$$

Introducing $\omega = e^{2L-1}$, the conditions for stability can be written

$$\begin{split} & 2(\omega-e^{-1})<1\\ & 2(\omega-e^{-1})>-1+(2(1-\omega)-e^{-1})\\ & 2(\omega-e^{-1})>-1-(2(1-\omega)-e^{-1}) \end{split}$$

From this follows that

$$\begin{aligned} &\omega < \frac{1}{2} + e^{-1} \\ &\omega > \frac{1 + e^{-1}}{4} \end{aligned}$$

The first inequality leads to

$$2L - 1 < log(1/2 + e^{-1}) = -0.1417$$

from which follows that

From the second inequality we have that

$$2L - 1 > log(1/4 + e^{-1}/4) = -1.0730$$

from which follows that

$$L > -0.0365$$

Hence, the largest delay L for which the system is stable is L = 0.4291.

4.1 Assuming the sampling interval h, use the various methods below to determine discrete-time approximations of the stable transfer function

$$G(s) = \frac{a}{s+a}, \qquad a > 0$$

For what values of h is the discrete-time system stable? For what values of h is the pole on the positive real axis? (What is the qualitative behavior of the discrete-time system if the pole is on the negative real axis?)

- a. Forward difference (Euler's method)
- b. Backward difference
- c. Tustin's method
- **d.** Tustin's method with prewarping and warping frequency $\omega_1 = a \text{ rad/s}$

Solution

a. Euler's method implies a translation

We get

$$s' = \frac{z-1}{h}$$

$$H(z) = G(s') = \frac{a}{\frac{z-1}{h} + a} = \frac{ah}{z + (ah-1)}$$

The discrete-time pole is located in z = 1 - ah. The system is stable if the pole is inside the unit circle. This is the case for $h < \frac{2}{a}$.

For $h < \frac{1}{a}$ the pole ends up on the positive real axis, and the discretetime system will have a somewhat similar behavior to the continuous-time system. For $h > \frac{1}{a}$, the pole will be on the negative real axis, and the system will oscillate – completely unlike the continuous system.

b. The backward difference implies a translation

$$s' = \frac{z-1}{zh}$$

We get

$$H(z)=G(s')=\frac{a}{\frac{z-1}{zh}+a}=\frac{zah}{z(1+ah)-1}$$

The discrete-time pole is located in $z = \frac{1}{1+ah}$. The system is stable and the pole is on the positive real axis for all values of h.

c. We use $s' = \frac{2}{h} \frac{z-1}{z+1}$ and obtain

$$H(z) = G(s') = \frac{a}{\frac{2}{h}\frac{z-1}{z+1} + a} = \frac{ah(z+1)}{2(z-1) + ah(z+1)} = \frac{ah(z+1)}{(2+ah)z + (ah-2)}$$

The discrete-time pole is located in $z = \frac{2-ah}{2+ah}$. The system is stable for all values of *h*, since

$$\left|\frac{2-ah}{2+ah}\right| < 1 \quad \Leftrightarrow \quad -2 < 2 < 2 + 2ah$$

which holds for all positive values of a and h. The pole ends up on the positive real axis for $h < \frac{2}{a}$.

d. We use $s' = \frac{a}{\tan(ah/2)} \frac{z-1}{z+1} = \frac{a}{\gamma} \frac{z-1}{z+1}$ and obtain

$$H(z) = G(s') = \frac{\alpha}{\frac{\alpha' z - 1}{\gamma z + 1} + \alpha'} = \frac{1}{\frac{1}{\gamma z - 1} + 1} = \frac{\gamma(z+1)}{\frac{\gamma(z+1)}{(z-1) + \gamma(z+1)}} = \frac{\gamma(z+1)}{(1+\gamma)z + \gamma - 1}$$

The discrete-time pole is thus located in

$$\bar{z} = \frac{1-\gamma}{1+\gamma} = \frac{1-\tan(ah/2)}{1+\tan(ah/2)} = \frac{1+1-(1+\tan(ah/2))}{1+\tan(ah/2)} = \frac{2}{\tan(ah/2)+1} - 1$$

which is a real number $\forall a \in \mathbb{R}, h > 0$. If we analyze the first term, it is easy to see that it goes from $-\infty$ to ∞ by changing a and/or h within their domains. If we limit the analysis for a while in the range $0 < ah < 3\pi/2$ (since it is the domain of \bar{z}), we can see that \bar{z} is monotonically decreasing

$$\frac{d\bar{z}}{d\,ah} = -\frac{\sec^2\left(\frac{ah}{2}\right)}{\left(\tan\left(\frac{ah}{2}\right)+1\right)^2} < 0$$

In addition, we know that for ah = 0 then $\bar{z} = 1$, and for $ah = \pi/2$ then $\bar{z} = -\infty$. The limit value for ah is then the one that gives $\bar{z} = -1$, i.e., when $ah = \pi$.

In other words, we can say that the discrete-time system is asymptotically stable whenever $0 < ah < \pi$. Since the value \bar{z} is a periodic function it is easy to extend the result as $2k\pi < ah < (2k+1)\pi$, with $k \in \mathbb{N}$ (recall that a > 0 and h > 0).

Figure 4.2 illustrates the plot of \bar{z} as a function of ah.



Figure 4.2 Plot of \bar{z} as a function of ah.

Given the analysis above, it is easy to say that the eigenvalue ends up on the positive real axis when $2k\pi < ah < 2k\pi + \pi/2$ with $k \in \mathbb{N}$.

6.1 Consider the task set below.

Task name	T_i	D_i	C_i
A	10	2	1
В	5	4	2
С	20	10	4

- **a.** Assign priorities to the tasks according to the rate-monotonic principle. Draw the schedule assuming worst-case conditions, i.e., the tasks are released simultaneously and the actual execution times are equal to the worst-case execution times. Do the tasks meet their deadlines?
- **b.** Assign priorities to the tasks according to the deadline-monotonic principle. Draw the schedule assuming worst-case conditions. Do the tasks meet their deadlines?

Solution

a. The tasks are given the following priorities:

Task name	T_i	Priority
А	10	Medium
В	5	High
С	20	Low

The schedule is shown in Figure 6.3. The worst-case response times of the tasks are $R_A = 3$, $R_B = 2$, $R_C = 9$, i.e., task A will not meet its deadlines.



Figure 6.3 Schedule with rate-monotonic priority assignments.

b. The tasks are given the following priorities:

Task name	D_i	Priority
А	2	High
В	4	Medium
С	10	Low

The schedule is shown in Figure 6.4. The worst-case response times of the tasks are $R_A = 1$, $R_B = 3$, $R_C = 9$, i.e. all tasks will meet their deadlines.



Figure 6.4 Schedule with deadline-monotonic priority assignments.